## Computer Science Tripos, Part II: Denotational Semantics Supervision 3 Daniel Sääw<sup>1</sup> (dks28)

- 1. Is PCF Turing-complete? What conclusions can we draw about the relationship between computability and continuity? Which result from a previous supervision do you need to use to prove that the Halting problem is unsolvable?
- 2. What sets are denoted by
  - (a)  $PCF_{nat \rightarrow nat}$
  - (b) **[bool**]
  - (c) **[**ℕ]
  - (d)  $[\![\langle x \mapsto \texttt{nat}, y \mapsto \texttt{bool}, z \mapsto \texttt{nat} \to \texttt{bool} \rangle]\!]$
- 3. Are terms and expressions the same thing in PCF?
- 4. State the full, formal version of Proposition 5.4.1. Carry out the proof indicated in the lecture notes.
- 5. 2016, Paper 9, Question 5
- 6. What slide is the central result of this course presented on? What assumptions does its proof make?
- 7. 2009, Paper 9, Question 6
- 8. Define  $\Omega_{\tau} := \operatorname{fix}(\operatorname{fn} x : \tau, x)$ . Show that  $\llbracket \Omega_{\tau} \rrbracket$  is the least element of the domain  $\llbracket \tau \rrbracket$ . Then deduce that  $\llbracket \Omega_{\tau \to \tau} \rrbracket = \llbracket \operatorname{fn} x : \tau, \Omega_{\tau} \rrbracket$ .

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9. The Fibonacci sequence can be computed in many ways, including the two Python functions below:

```
def fib1(n):
    if n == 0:
        return 0
    elif n == 1:
        return 1
    else:
        return fib1(n-1) + fib1(n-2)
def fib2(n):
    a = 1
    b = 0
    while n > 0:
        a, b = a+b, a
        n -= 1
    return b
```

Define two PCF expressions computing the nth Fibonacci number in the two ways above. Show that they are contextually equivalent.

- 10. Exercise 7.4.1
- 11. Exercise 7.4.2
- 12. For each of the following statements, decide whether it is true or false, and justify why.
  - (a)  $\langle x \mapsto \texttt{nat}, y \mapsto \texttt{nat} \rangle \vdash \texttt{fn } y : \texttt{nat. succ}(y) \simeq_{\texttt{ctx}} \texttt{fn } x : \texttt{nat. succ}(x) : \texttt{nat} \to \texttt{nat}$
  - (b)  $x \simeq_{\text{ctx}} x$
  - (c) If  $V \in \text{PCF}_{\tau}$  is a value,  $M \in \text{PCF}_{\tau}$  and  $V \simeq_{\text{ctx}} M$ , then  $M \Downarrow_{\tau} V$ .
  - (d)  $\perp_{\mathbb{N}} \triangleleft_{\texttt{nat}} \texttt{if } \texttt{zero}(\texttt{succ}(\texttt{0})) \texttt{ then } x \texttt{ else } 0$
  - (e) Since extending PCF with por is sufficient to get full abstraction, por is the only function in  $\mathbb{B}_{\perp} \to \mathbb{B}_{\perp} \to \mathbb{B}_{\perp}$  that is not implementable in pure PCF.
- 13. 2021, Paper 9, Question 7

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